

Sisyphus: Mostly-Automated Proof Repair for Verified Libraries



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Pictured: OCaml programmer fixing a broken Coq proof

Let's write a program!

Let's write a **verified** program!

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Q: Convert a sequence to an array.

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$\forall s \ell, \{s \mapsto \text{Seq } \ell\}$

(`to_array` s)

$\exists a, \{a \mapsto \text{Array } \ell\}$

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OCaml's `'a Seq.t` datatype

```
type 'a t    = unit -> 'a node
and  'a node = Nil
      | Cons of 'a * (unit -> 'a node)
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A *thunked tail*

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```



```
fun () -> Cons (1, fun () -> Cons (2, fun () -> Nil))
```

Q: Convert a sequence to an array.

$$\forall s \ell, \{s \mapsto \text{Seq } \ell\}$$
$$(\text{to_array } s)$$
$$\exists a, \{a \mapsto \text{Array } \ell\}$$

Q: Convert a sequence to an array.

*finite and
effect-free*

$\forall s \ell, \{s \mapsto \text{Seq } \ell\}$

(to_array s)

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“ a ” **points-to** an array

Q: Convert a sequence to an array.

$\forall s \ell, \{s \mapsto \text{Seq } \ell\}$

(`to_array` s)

$\exists a, \{a \mapsto \text{Array } \ell\}$

Let's write
some
code/proofs!

$\forall s \ell, \{s \mapsto \mathbf{Seq} \ell\}$ $(\mathbf{to_array} s)$ $\exists a, \{a \mapsto \mathbf{Array} \ell\}$

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let to_array s = xcf.
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$\{s \mapsto \text{Seq } \ell\}$

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let to_array s =           xcf.  
  match s () with
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```
let to_array s =  
  match s () with
```

```
    xcf.  
    xapp; case  $\ell$  as [| h tl|]
```

$\{s \mapsto \text{Seq } []\}$

```
let to_array s = xcf.  
  match s () with xapp; case  $\ell$  as [|  $h\ t|]$   
  | Nil ->
```

$\{s \mapsto \text{Seq } []\}$

```
let to_array s =  
  match s () with  
  | Nil -> []
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xcf.
xapp; case ℓ as [| $h\ t|]$

$\{s \mapsto \text{Seq } []\}$

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let to_array s =  
  match s () with  
  | Nil -> []
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xcf.
xapp; case ℓ as [| $h\ t|]$
– xvalemptyarr.

$$\exists a, \{a \mapsto \text{Array } []\}$$

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let to_array s = xcf.  
  match s () with xapp; case l as [| h tl]  
  | Nil -> [] - xvalemptyarr.  
  | Cons (h, _) -> -  
    let sz = length s in
```

$$\{sz = \text{length } \ell; s \mapsto \text{Seq } (h :: tl)\}$$

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let to_array s =
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    let a = make sz h in
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xapp; case ℓ as [| $h\ tl$]
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$\{sz = \text{length } \ell; s \mapsto \text{Seq } (h :: tl) * a \mapsto \text{Array } (\text{make } sz \ h)\}$

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xapp; case ℓ as [| $h \ tl$]
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xapp.
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let to_array s =

xcf.

$t \text{ ++ } drop(\text{length } t) (\text{make } (\text{length } \ell) \ h)$

| **Cons** (h, _) ->

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let sz = length s **in**

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let to_array s = xcf.  
  t ++ drop (length t) (make (length ℓ) h)  
  | Cons (h, _) ->  
    let sz = length s in xapp.  
    let a = make sz h in xalloc.  
    iteri (fun i vl -> xapp (iteri_spec (λt →  
      a ↪ Array (  
        a.(i) <- vl t ++ drop (length t)  
        (make (length ℓ) h))  
      ) s;  
    ).
```

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let to_array s =

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Qed.

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Let's write a **verified** program!

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Conclusion:

Writing verified code is hard

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A problem arises..

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Make Seq.to_array behave better with stateful sequences #390

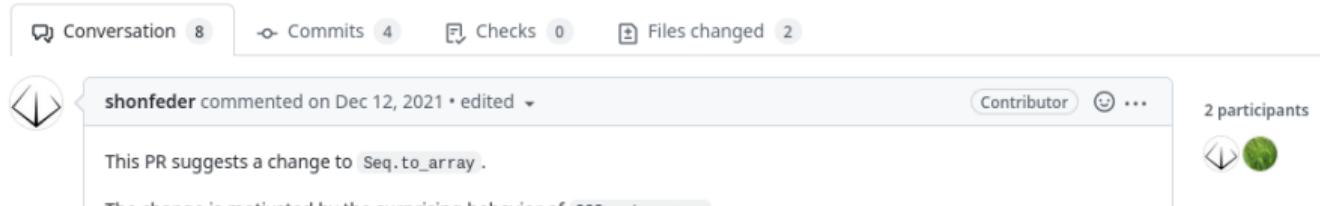
Merged c-cube merged 4 commits into [c-cube:master](#) from [shonfeder:state-friendly-seq-to-array](#) on Dec 12, 2021

Conversation 8 Commits 4 Checks 0 Files changed 2

shonfeder commented on Dec 12, 2021 • edited

This PR suggests a change to `Seq.to_array`.

Contributor ... 2 participants



Old

```
let to_array s =
  match s () with
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    let sz = length s in
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```

New

```
let to_array s =
  let sz, ls = fold_left
    (fun (i, acc) x ->
      (i+1, x::acc))
    (0, [])
  in
  match ls with
  | [] -> []
  | init :: rest ->
    let a = make sz init in
    let idx = len - 2 in
    List.fold_left
      (fun i vl ->
        a.(i)<-vl; i-1)
      idx
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    (fun (i, acc) x ->
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    (0, []) l in
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Completely different implementation...

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```

Completely different implementation...

...proof must be redone.

Writing verified code is hard...

Maintaining
~~Writing~~ verified code is hard...

Maintaining

~~Writing~~ verified code is hard...

...can we make it easier?

Contributions

- 1 **Sisyphus**: tool to *repair* proofs over changes.
- 2 **PDT**: technique to *efficiently* test invariants.
- 3 Evaluation on 10 real OCaml programs and their changes

New Program:

```
let to_array s =  
  
  let sz, ls = fold_left  
    (fun (i, acc) x ->  
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      idx rest;  
    a
```

New Program:

```
let to_array s =  
  
    let sz, ls = fold_left  
        ~f:(fun a l -> a :: l)  
        ~init:[], s  
    in  
    match ls with  
    | [] -> [] []  
    | init :: rest ->  
        let a = make sz init in  
        let idx = sz - 2 in  
        List.fold_left  
            (fun i vl ->  
                a.(i)<-vl; i-1)  
            idx rest;  
        a
```

How to generate a proof script?

New Program:

```
let to_array s =  
  
    let sz, ls = fold_left
```

How to generate a proof script?

```
        ~> ls :: s  
match ls with  
| [] -> [] []  
| init :: rest ->
```

Observation: proofs are syntax-directed

```
List.fold_left  
  (fun i vl ->  
    a.(i)<-vl; i-1)  
  idx rest;  
  a
```

New Program:

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let to_array s =  
  
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How to generate a proof script?

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        ~> ls :: s  
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How to fill in the holes?

Key challenges for proof repair

- 1 Generating candidate invariants
- 2 Testing generated invariants

Key challenges for proof repair

1 Generating candidate invariants

2 Testing generated invariants

Generating candidate invariants

How to fill in holes?

Generating candidate invariants

How to fill in holes?

Use the old program and proofs...

Generating candidate invariants

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let to_array l =
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Generating candidate invariants

Programs are different...

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...but have similarities.

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Generating candidate invariants

Q: How to discover these similarities automatically?

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A: Instrumentation based dynamic analysis

Generating candidate invariants

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```

A: Instrumentation based dynamic analysis

Identify “similar” program points through traces.

Generating candidate invariants

Q: How to discover these similarities automatically?

```
let to_array l =  
  match l () with  
  | Nil ->  
  | Cons (x, xs) ->  
    let len = length' l in  
    let rec iteri i a =  
      (fun x -> a.(i) <- x)  
    in iteri (len - 1) l
```

```
let to_array l =  
  let sz, ls =  
    fold (fun (i, acc) x ->  
          (0, [ ]) l in  
    match ls with  
    | [] -> []  
    | init::rest ->  
      let idx = sz - 2 in  
      let _ =  
        List.fold_left  
        (fun i x -> a.(i) <- x; i + 1)  
        a
```

A: Instrumentation based dynamic analysis

Identify “similar” program points through traces.

Use invariants from old proof to synthesise invariants for new one.

Key challenges for proof repair

- 1 *Generating candidate invariants*

- 2 *Testing generated invariants*

Key challenges for proof repair

- 1 *Generating candidate invariants*
- 2 *Testing generated invariants*

Testing Candidate Invariants

Testing Candidate Invariants

```
xcf.  
  
xapp (...).  
  
  case l as [] init rest].  
    - xmatch 0. xvaleemptyarr.  
    - xmatch 1.  
      xalloc a data Hdata.  
      xlet idx.  
      xapp (fold_left_spec idx rest  
        (fun acc ls =>  
          (??))).  
  
xvals.
```

Testing Candidate Invariants

```
xcf.  
  
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              (fun acc ls =>  
                (??))).  
  
xvals.
```

Testing Candidate Invariants

fold_left_spec

Testing Candidate Invariants

fold_left_spec:

$$\forall I f acc \ell, \left(\begin{array}{c} \{I acc' t\} \\ \forall acc' v t, \quad (f acc' v) \\ \exists res, \{I res (t ++ [v])\} \end{array} \right) \rightarrow$$

$$\begin{aligned} & \{I acc []\} \\ & (\text{List}. \text{fold_left } f acc \ell) \\ & \exists res, \{I res \ell\} \end{aligned}$$

Testing Candidate Invariants

fold_left_spec:

Loop Invariant

$$\forall I f \text{ acc } \ell, \left(\forall \text{acc}' v t, \begin{array}{c} \{I \text{ acc}' t\} \\ (f \text{ acc}' v) \\ \exists \text{res}, \{I \text{ res} (t \text{ ++ } [v])\} \end{array} \right) \rightarrow$$

$$\begin{array}{c} \{I \text{ acc } []\} \\ (\text{List. fold_left } f \text{ acc } \ell) \\ \exists \text{res}, \{I \text{ res } \ell\} \end{array}$$

Testing Candidate Invariants

fold_left_spec:

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$$\begin{array}{c} \{I acc []\} \\ (\text{List}. \text{fold_left } f acc \ell) \\ \exists res, \{I res \ell\} \end{array}$$

*Preservation
of invariant*

Testing Candidate Invariants

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Testing Candidate Invariants

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Testing Candidate Invariants

fold_left_spec:

If I holds on t

$$\forall I f acc \ell, \left(\begin{array}{c} \forall acc' v t, \quad \{I acc' t\} \\ \quad (f acc' v) \\ \exists res, \{I res (t ++ [v])\} \end{array} \right) \rightarrow$$

$$\begin{aligned} & \{I acc []\} \\ (\text{List}. & \text{fold_left } f acc \ell) \\ & \exists res, \{I res \ell\} \end{aligned}$$

Testing Candidate Invariants

fold_left_spec:

If I holds on t

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after calling f...

$$\begin{aligned} & \{I acc []\} \\ & (\text{List}. \text{fold_left } f acc \ell) \\ & \exists res, \{I res \ell\} \end{aligned}$$

Testing Candidate Invariants

fold_left_spec:

If I holds on t

$$\forall I f acc \ell, \left(\begin{array}{c} \forall acc' v t, \quad \{I acc' t\} \\ \exists res, \{I res (t ++ [v])\} \\ (f acc' v) \end{array} \right) \rightarrow$$

after calling f...

$$\begin{aligned} & \{I acc []\} \\ & (\text{List}. \text{fold_left } f acc \ell) \\ & \exists res, \{I res \ell\} \end{aligned}$$

I holds on t ++ [v]

Testing Candidate Invariants

fold_left_spec:

$$\forall I f acc \ell, \left(\begin{array}{c} \{I acc' t\} \\ \forall acc' v t, \quad (f acc' v) \\ \exists res, \{I res (t ++ [v])\} \end{array} \right) \rightarrow$$

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Testing Candidate Invariants

fold_left_spec:

$$\forall I f acc \ell, \left(\begin{array}{c} \{I acc' t\} \\ \forall acc' v t, \quad (f acc' v) \\ \exists res, \{I res (t ++ [v])\} \end{array} \right) \rightarrow$$

Initial condition $\{I acc []\}$
 $(List.fold_left f acc \ell)$
 $\exists res, \{I res \ell\}$

Testing Candidate Invariants

fold_left_spec:

$$\forall I f acc \ell, \left(\begin{array}{c} \{I acc' t\} \\ \forall acc' v t, \quad (f acc' v) \\ \exists res, \{I res (t ++ [v])\} \end{array} \right) \rightarrow$$

$\{I acc []\}$
 $(List.\text{fold_left } f acc \ell)$
 $\exists res, \{I res \ell\}$ *I holds on entire ℓ*

Testing Candidate Invariants

fold_left_spec:

How to test a candidate for I ?

$$\begin{array}{c} \{I\ acc\ []\} \\ (\text{List}. \text{fold_left } f\ acc\ \ell) \\ \exists res, \{I\ res\ \ell\} \end{array}$$

Testing Candidate Invariants

fold_left_spec:

$$\forall I f acc \ell, \left(\begin{array}{c} \{I acc' t\} \\ \forall acc' v t, (f acc' v) \\ \exists res, \{I res (t ++ [v])\} \end{array} \right) \rightarrow$$

$\{I acc []\}$
 $(List.fold_left f acc \ell)$
 $\exists res, \{I res \ell\}$

Depends on
logical variables

Testing Candidate Invariants

fold_left_spec:

How to test a candidate for I ?

$$\begin{array}{c} \{I\ acc\ []\} \\ (\text{List}. \text{fold_left } f\ acc\ \ell) \\ \exists res, \{I\ res\ \ell\} \end{array}$$

Testing Candidate Invariants

`fold_left_spec:`

How to test a candidate for I ?

$$\forall I f acc \ell, \left(\begin{array}{c} \forall acc' v t, \\ \exists res, \{I res (t ++ [v])\} \end{array} \right) \rightarrow$$

$$\begin{aligned} & \{I acc []\} \\ & (\text{List}.fold_left f acc \ell) \\ & \quad \exists res, \{I res \ell\} \end{aligned}$$

Testing Candidate Invariants

fold_left_spec:

How to test a candidate for I ?

$$\forall I f \text{ acc } \ell, \left(\forall \text{acc}' v t, \quad \frac{}{\text{acc}' v} \right) \rightarrow$$

Idea: Use the proof term of fold_left_spec

$$\begin{aligned} & \{I \text{ acc } []\} \\ & (\text{List}. \text{fold_left } f \text{ acc } \ell) \\ & \exists \text{res}, \{I \text{ res } \ell\} \end{aligned}$$

Testing Candidate Invariants

```
let fold_left f acc ls =  
  match ls with  
  | [] ->  
    acc  
  | h :: t ->  
    let acc' = f acc h in  
    fold_left f acc' t
```

Testing Candidate Invariants

```
let fold_left f acc ls =
  {I acc ℓ'}
  match ls with
  | [] ->
    acc
  | h :: t ->
    let acc' = f acc h in
    fold_left f acc' t
```

Testing Candidate Invariants

```
let fold_left f acc ls =
  {I acc l'}
  match ls with
  | [] ->
    {I acc l'}
    acc
  | h :: t ->

    let acc' = f acc h in
    fold_left f acc' t
```

Testing Candidate Invariants

```
let fold_left f acc ls =
  {I acc l'}
  match ls with
  | [] ->
    {I acc l'}
    acc
  | h :: t ->
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    let acc' = f acc h in
      fold_left f acc' t
```

Testing Candidate Invariants

```
let fold_left f acc ls =
  {I acc ℓ'}
  match ls with
  | [] ->
    {I acc ℓ'}
    acc
  | h :: t ->
    {I acc ℓ'}
    let acc' = f acc h in
    {I acc' (ℓ' ++ [h])}
    fold_left f acc' t
```

Testing Candidate Invariants

```
let fold_left f acc ls =
  {I acc ℓ'}
  match ls with
  | [] ->
    {I acc ℓ'}
    acc
  | h :: t ->
    {I acc ℓ'}
    let acc' = f acc h in
    {I acc' (ℓ' ++ [h])}
    fold_left f acc' t
    {I acc (ℓ' ++ ℓ)}
```

Testing Candidate Invariants

```
let fold_left f acc ls =  
  {I acc ℓ'}  
  match ls with  
  | [] ->  
    {I acc ℓ'}  
    acc  
  | h :: t ->  
    {I acc ℓ'}  
    let acc' = f acc h in  
    {I acc' (ℓ' ++ [h])}  
    fold_left f acc' t  
    {I acc (ℓ' ++ ℓ)}
```



Describes **exactly**
how *I* is maintained

Testing Candidate Invariants

```
let fold_left f acc ls =  
  {I acc ℓ'}  
  match ls with  
  | [] ->  
    {I acc ℓ'}  
    acc  
  | h :: t ->  
    {I acc ℓ'}  
    let acc' = f acc h in  
    {I acc' (ℓ' ++ [h])}  
    fold_left f acc' t  
    {I acc (ℓ' ++ ℓ)}
```



Describes **exactly**
how *I* is maintained

Proof-driven Testing

Testing Candidate Invariants

fold_left_spec

Testing Candidate Invariants

```
fold_left_spec ?I f 2 [2;1] ?HI
```

Testing Candidate Invariants

Instantiate with concrete arguments..

```
fold_left_spec ?I f 2 [2;1] ?HI
```

Testing Candidate Invariants

Instantiate with concrete arguments..

```
fold_left_spec ?I f 2 [2;1] ?HI
```

Testing Candidate Invariants

Instantiate with concrete arguments..

```
fold_left_spec ?I f 2 [2;1] ?HI
```

...with existentials for proof arguments

Testing Candidate Invariants

```
fold_left_spec ?I f 2 [2;1] ?HI
```

Testing Candidate Invariants

```
fold_left_spec ?I f 2 [2;1] ?HI
```

Testing Candidate Invariants

```
fold_left_spec ?I f 2 [2;1] ?H
```

|

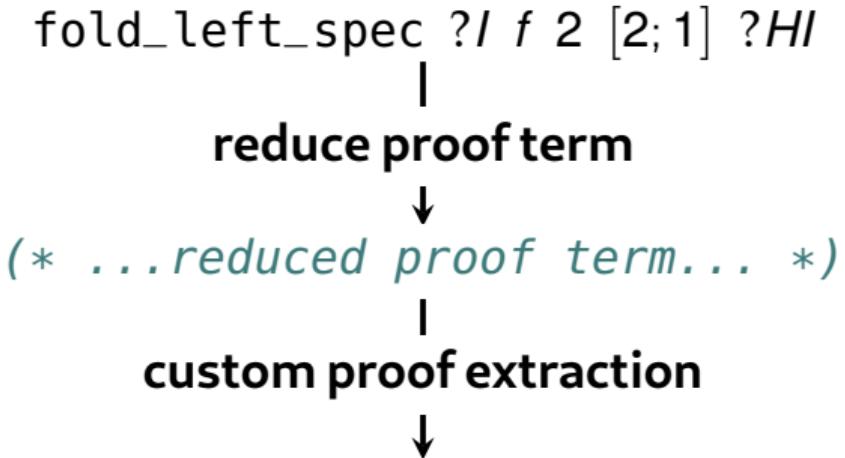
reduce proof term

↓

Testing Candidate Invariants

```
fold_left_spec ?I f 2 [2;1] ?HI  
|  
reduce proof term  
↓  
(* ...reduced proof term... *)
```

Testing Candidate Invariants



Testing Candidate Invariants

```
assert (l len []);
let acc = f len 2 in
assert (l acc [2]);
let acc = f acc 1 in
assert (l acc [2; 1]);
()
```

Testing Candidate Invariants

```
assert (l len []);
let acc = f len 2 in
assert (l acc [2]);
let acc = f acc 1 in
assert (l acc [2; 1]);
()
```

Simulates an execution of `List.fold_left`

Testing Candidate Invariants

Instantiate I with embedding of candidate invariant...

```
assert (I len []);
let acc = f len 2 in
assert (I acc [2]);
let acc = f acc 1 in
assert (I acc [2; 1]);
()
```

Testing Candidate Invariants

Instantiate I with embedding of candidate invariant...

```
assert (I len []);
let acc = f len 2 in
assert (I acc [2]);
let acc = f acc 1 in
assert (I acc [2; 1]);
()
```

...prune candidate if assertion raised.

Key challenges for proof repair

- 1 *Generating candidate invariants*
- 2 *Testing generated invariants*

Key challenges for proof repair

- 1 Generating candidate invariants
...using *relevant* expressions from old proof

- 2 Testing generated invariants

Key challenges for proof repair

- 1 Generating candidate invariants
...using *relevant* expressions from old proof

- 2 Testing generated invariants
...using *proof-driven testing*

Research Questions

- 1 Is Sisyphus effective at repairing proofs?
- 2 Does Sisyphus repair proofs in reasonable time?
- 3 What changes does Sisyphus handle poorly?

Benchmark Programs

- 14 OCaml programs and their changes
- 10 from real-world OCaml codebases
 - ...such as containers or Jane Street's core

Benchmark Programs

Example	Data Structure	Refactoring
seq_to_array	Array, Seq	IterOrd, DataStr
make_rev_list [†]	Ref	Mutable/Pure
tree_to_array [†]	Array, Tree	IterOrd, DataStr
array_exists	Array	Mutable/Pure
array_find_mapi	Array, Ref	Pure/Mutable
array_is_sorted	Array	Pure/Mutable
array_findi	Array	Pure/Mutable
array_of_rev_list	Array	DataStr
array_foldi	Array	Pure/Mutable
array_partition	Array	DataStr
stack_filter [‡]	Stack	DataStr
stack_reverse [‡]	Stack	DataStr
sll_partition [‡]	SLL	Mutable/Pure, IterOrd
sll_of_array [‡]	Array, SLL	IterOrd

RQ1: Effectiveness of proof repair

Name	Time (old)			# Admits / # Obligations	Time (new)
	Spec	Proof	Total		
seq_to_array	1hrs	1hr	2hrs	3 / 5	17m
make_rev_list	5m	5m	10m	0 / 2	-
tree_to_array	4hrs	1hr	5hrs	2 / 4	18m
array_exists	10m	20m	30m	2 / 4	12m
array_find_mapi	30m	1hr	1.5hrs	2 / 5	12m
array_is_sorted	1hr	3hrs	4hrs	2 / 5	2m
array_findi	30m	1hr	1.5hrs	3 / 7	9m
array_of_rev_list	5m	1hr	1hr	2 / 3	3m
array_foldi	10m	5m	15m	0 / 1	-
array_partition	30m	2hrs	2.5hrs	3 / 3	5m
stack_filter	1hr	30m	1.5hrs	3 / 3	11m
stack_reverse	1.5hrs	30m	2hrs	1 / 1	30s
sll_partition	1hr	1hr	2hrs	0 / 2	-
sll_of_array	1.5hrs	30m	2hrs	0 / 1	-

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array_partition	30m	2hrs	2.5hrs	3 / 3	5m
stack_filter	1hr	30m	1.5hrs	3 / 3	11m
stack_reverse	1.5hrs	30m	2hrs	1 / 1	30s
sll_partition	1hr	1hr	2hrs	0 / 2	-
sll_of_array	1.5hrs	30m	2hrs	0 / 1	-

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array_of_rev_list	5m	1hr	1hr	2 / 3	3m
array_foldi	10m	5m	15m	0 / 1	-
array_partition	30m	2hrs	2.5hrs	3 / 3	5m
stack_filter	1hr	30m	1.5hrs	3 / 3	11m
stack_reverse	1.5hrs	30m	2hrs	1 / 1	30s
sll_partition	1hr	1hr	2hrs	0 / 2	-
sll_of_array	1.5hrs	30m	2hrs	0 / 1	-

RQ2: Efficiency of proof repair

Example	Time (s)				Total (s)
	Generation	Extraction	Testing	Remaining	
seq_to_array	28.57	1.95	20.36	5.28	58
make_rev_list	$\leq 10ms$	3.36	$\leq 10ms$	11.95	15
tree_to_array	6.75	1.95	2.98	13.32	25
array_exists	$\leq 10ms$	3.30	$\leq 10ms$	13.23	17
array_find_mapi	$\leq 10ms$	2.13	$\leq 10ms$	13.95	17
array_is_sorted	$\leq 10ms$	2.04	$\leq 10ms$	15.38	18
array_findi	$\leq 10ms$	2.13	$\leq 10ms$	19.07	22
array_of_rev_list	1.72	2.82	0.96	15.62	21
array_foldi	$\leq 10ms$	488.89	$\leq 10ms$	15.00	504
array_partition	3.51	69.73	2.62	17.53	95
stack_filter	$\leq 10ms$	81.88	$\leq 10ms$	21.53	104
stack_reverse	$\leq 10ms$	88.42	$\leq 10ms$	16.94	105
sll_partition	$\leq 10ms$	426.62	$\leq 10ms$	16.43	443
sll_of_array	0.02	55.98	0.01	13.33	69

RQ2: Efficiency of proof repair

Example	Time (s)				Total (s)
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stack_filter	$\leq 10ms$	81.88	$\leq 10ms$	21.53	104
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sll_partition	$\leq 10ms$	426.62	$\leq 10ms$	16.43	443
sll_of_array	0.02	55.98	0.01	13.33	69

RQ3: Failure Modes

-  Repair assumes components from old proof are sufficient for new one.
-  Quality of repair degrades when this fails to hold.

e.g. `array_partition`'s pure obligations required fact

$$\text{filter } p(\text{filter } p \ell) = \text{filter } p \ell$$

not present in original proof.

RQ3: Failure Modes

- Repair assumes components from old proof are sufficient for new one.
- Quality of repair degrades when this fails to hold.
- In some cases, repair may even fail.

RQ3: Failure Modes

```
let to_array s =
  let batches = (* .. *) in
  let res =
    Array.make (* .. *) in
  List.iter (fun batch ->
    let dst = (* .. *) in
    Array.copy batch res dst)
  batches;
res
```

RQ3: Failure Modes

Invariant requires *flattening* operation...

```
let to_array s =
  let batches = (* .. *) in
  let res =
    Array.make (* .. *) in
  List.iter (fun batch ->
    let dst = (* .. *) in
    Array.copy batch res dst)
  batches;
res
```

RQ3: Failure Modes

Invariant requires *flattening* operation...

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let to_array s =
  let batches = (* .. *) in
  let res =
    Array.make (* .. *) in
  List.iter (fun batch ->
    let dst = (* .. *) in
    Array.copy batch res dst)
  batches;
res
```

...not present in old proof.

Summary

- 1 **Sisyphus**: tool to *repair* proofs over changes.
- 2 **PDT**: technique to *efficiently* test invariants.
- 3 Evaluation on 10 real OCaml programs and their changes